

A language over a one symbol alphabet
 requiring only $O(\log \log n)$ space

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It is well known that the minimal growth function for the tape complexity of Turing machines is $\log \log n$ [2]. In the literature, one can find essentially one example of a language requiring only $O(\log \log n)$ space, namely

$$L_0 = \{ \# \text{bin}(1) \# \text{bin}(2) \# \text{bin}(3) \# \dots \# \text{bin}(n) \# ; n \in \mathbb{N} \}$$

where $\text{bin}(i)$ is the binary representation of the integer i .

In this note we describe a language over a one symbol alphabet having space complexity $O(\log \log n)$.

Let $L_1 = \{ a^n ; \text{the smallest number } q \text{ which does not divide } n \text{ is a power of two} \}$

For every natural number n let $q(n)$ be the smallest number which does not divide n .

Lemma: $\exists c > 0: q(n) \leq c \log n$

Proof: Let n be any natural number and let $m = q(n)$. Then all primes less than m divide n and hence their product P divides them. By the prime number theorem of Gauss there are about $\frac{m}{\log m}$ primes less than m and hence P is larger than

$(\frac{m}{\log m})! \geq 2^{m/4}$ for sufficiently large m .

q.e.d.

As an immediate consequence of lemma 1 we obtain that the following recognition procedure has space complexity $O(\log \log n)$.

input $n \in \mathbb{N}$;

$q \leftarrow 1$;

repeat $q \leftarrow q + 1$; $p \leftarrow n \bmod q$ until $p \neq 0$;

if q is a power of two then accept else reject;

It remains to be shown that L_1 is not regular. For $k \in \mathbb{N}$ let m_k be the least common multiple of the numbers $1, \dots, 2^k - 1$. Then $a^{m_k} \in L_1$ for all k . Furthermore, since for infinitely many k there exists a prime between 2^k and 2^{k+1} , $a^{2m_k} \notin L_1$ for infinitely many k . Assume now that L is regular.

Then $L_1 = \bigcup_{i \in J} s_i + t_i \mathbb{N}$ for some finite set J and some $s_i, t_i \in \mathbb{N}_0$. For sufficiently large k every non-zero t_i divides m_k and hence with $m_k \in s_i + n_i \mathbb{N}$ also $2m_k \in s_i + t_i \mathbb{N}$. Contradiction!

Conclusion: We have exhibited a non-regular language over a one symbol alphabet requiring only $O(\log \log n)$ space to recognize. This example was found in the futile attempt to extend the result from [1]. There we showed:

Theorem: Let L be context-free and let u, v, w, x, y be words with $L \cap uv^*wx^*y$ non-regular. Then $\log n$ is a lower bound for the space complexity of L .

- [1] H. Alt & K. Mehlhorn: Untere Schranken für den Platzbedarf bei der kontext-freien Analyse, Techn. Bericht, Fachbereich 10 der Universität des Saarlandes, 1975
- [2] J. Hopcroft & J. Ullman: Formal Languages and their Relation to Automata, Addison-Wesley, Reading, Mass., 1969

How to Gracefully Number
Certain Symmetric Trees

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Let G be an (undirected) graph, with vertex set V and edge set E . A vertex-numbering assigns a distinct non-negative integer to each vertex of G . We regard a vertex and its vertex number as identical objects; thus V is the set of all vertex numbers. With an edge (v, w) we associate the edge number $|v-w|$. We identify E with the set of all edge numbers. A graceful numbering of a tree with n vertices has $V = \{i \mid 0 \leq i < n\}$, $E = \{i \mid 1 \leq i < n\}$. For example, the trees in Figure 1(a) and 1(c) are gracefully numbered.

In $[C, R]$, constructions are given for gracefully numbering a restricted class of symmetric trees. We give a more general construction, that gracefully numbers a tree if the subtrees of the root are all isomorphic to one gracefully numbered tree. We indicate other trees that can be gracefully numbered by our technique.